## **Advanced Robotics**

## Lecture 8

Algebraic equations & affine varieties

aly equations  $f_{k}(x_{1}, x_{2}, \dots, x_{m}) = 0$   $k = 1, \dots, 5$ polynomial  $f_{k}(x_{1}, x_{2}, \dots, x_{m}) = \sum_{\alpha} a_{\alpha} x^{\alpha}$ monomial  $X^{\alpha} \equiv X_1^{\alpha_1} \cdot X_2^{\alpha_2} \cdots \times_n^{\alpha_n}$  $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_m)$  $d_i \in \mathbb{Z}_{\geq 0}$ nonnegative whole numbers  $a_{\alpha} \in k \dots a field (R, C, \dots)$  $k[x_1, x_2, ..., x_m] = \text{the set of all polynomials in } x_1, x_2, ..., x_m$ (X1,X2,...,Xm) E k = n-dimensional linear space

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$$\begin{array}{c} \operatorname{Ring} af \ polynomials \qquad \text{like } \operatorname{R} \ \text{mithout the detrivion} \\ \end{array}$$

$$\begin{array}{c} \operatorname{Plus} \oplus \\ \operatorname{plus} \oplus \\ \operatorname{k}[x_{i_1}x_{i_1\cdots i_n}x_m] \times k[x_{i_1}x_{2}\cdots i_m] \longrightarrow k[x_{i_1}x_{i_1\cdots i_n}x_m] \\ \end{array}$$

$$\begin{array}{c} \operatorname{resc} \operatorname{fog} g = h \quad \text{or} \quad \operatorname{fog} g = m \\ \end{array}$$

$$\begin{array}{c} \operatorname{Resc} f \oplus g = h \quad \text{or} \quad \operatorname{fog} g = m \\ \end{array}$$

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$$\begin{array}{c} \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} x^{\alpha} = \sum_{\substack{\alpha \in A}} a_{\alpha}' x^{\alpha} \quad \text{where} \quad a_{\alpha}' = \left( \begin{array}{c} a_{\alpha} & x \in A \\ 0 & x \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha}' x^{\alpha} \quad \text{where} \quad a_{\alpha}' = \left( \begin{array}{c} a_{\alpha} & x \in A \\ 0 & x \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha}' x^{\alpha} \quad \text{where} \quad a_{\alpha}' = \left( \begin{array}{c} a_{\alpha} & y \in A \\ 0 & x \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha}' x^{\alpha} \quad \text{where} \quad b_{\alpha}' = \left( \begin{array}{c} b_{\alpha} & y \in A \\ 0 & y \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha}' x^{\alpha} \quad \text{where} \quad b_{\alpha}' = \left( \begin{array}{c} b_{\alpha} & y \in A \\ 0 & y \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha}' x^{\alpha} \quad \text{where} \quad b_{\alpha}' = \left( \begin{array}{c} b_{\alpha} & y \in A \\ 0 & y \in A \end{array} \right) \\ \operatorname{Resc} g = \sum_{\substack{\alpha \in A}} b_{\alpha} \left( a_{\alpha}' a_{\alpha}' x^{\alpha} \right) \\ \operatorname{Resc} x^{\alpha} x^{\alpha} = \sum_{\substack{\alpha \in A}} b_{\alpha} \left( \begin{array}{c} b_{\alpha} & a_{\alpha}' x^{\alpha} \\ a_{\alpha} & a_{\alpha}' x^{\alpha} & a_{\alpha}' x^{\alpha} \end{array} \right) \\ \operatorname{Resc} f \oplus g = \sum_{\substack{\alpha \in A \cup B}} \sum_{\substack{\alpha \in A \cup B}} (a_{\alpha}' a_{\alpha}' x^{\alpha}) \\ \operatorname{Resc} x^{\alpha} x^{\alpha} = \sum_{\substack{\alpha \in A \cup B}} \sum_{\substack{\alpha \in A \cup B}} (a_{\alpha}' a_{\alpha}' x^{\alpha}) \\ \operatorname{Resc} x^{\alpha} x^{$$

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Excercises:

Show that (k[x<sub>1</sub>,...,k<sub>n</sub>],⊕,⊙) is a ring, i.e. that holds true
 (a) f ⊕ g = g ⊕ f

(b) 
$$(f \oplus g) \oplus h = f \oplus (g \oplus h)$$

- (c)  $\exists 0 \in k[x_1,\ldots,k_n]$ :  $f \oplus 0 = f$
- (d)  $\forall f \exists e \in k[x_1, \ldots, k_n]: e \oplus f = 0$

(e) 
$$f \odot g = g \odot f$$

(f) 
$$(f \odot g) \odot h = f \odot (g \odot h)$$

(g) 
$$\exists 1 \in k[x_1,\ldots,k_n]$$
:  $f \odot 1 = f$ 

(h) 
$$f \odot (g \oplus h) = (f \odot g) \oplus (f \odot h)$$

for arbitrary  $f, g, h \in k[x_1, \ldots, k_n]$ .

$$f = \sum_{i} a_{i} x^{i} x^{i} = a_{(0,0,\dots,0)} x^{(0,0,\dots,0)} + \dots + a_{(\alpha_{i},\alpha_{2},\dots,\alpha_{m})} x^{(\alpha_{i},\alpha_{2},\dots,\alpha_{m})} + \dots$$

$$f = S(k_{i},k_{2},\dots,k_{m}) \quad S(k_{i},k_{2},\dots,k_{m}) = \begin{pmatrix} a_{i} & \text{if } \exists a \in A : (k_{i},k_{2},\dots,k_{m}) = a \\ 0 & \text{otherwise} \end{pmatrix}$$

$$f = S(k_{i},k_{2},\dots,k_{m}) \quad S(k_{i},k_{2},\dots,k_{m}) = \begin{pmatrix} a_{i} & \text{if } \exists a \in A : (k_{i},k_{2},\dots,k_{m}) = a \\ 0 & \text{otherwise} \end{pmatrix}$$

$$f(x) = 2x^{3} + x + 3 = 3x^{0} + 1 \cdot x' + 0 \cdot x^{2} + 2x^{3} + 0 \cdot x^{4} + 0 \cdot x^{5} + \dots$$

$$= (3 + 1 + 0 + 2 + 0 + 0 + 1 + 0) \dots \text{ on infinite}$$

$$x^{0} \times^{4} \times^{2} \times^{3} \times^{4} \times^{5} \dots \qquad \text{sequence}$$

$$\text{ with non-zero values}$$

$$\text{ on a finite number}$$

$$\text{ of coordinates}$$

$$\text{ but we often write } (2 + 0 + 1 + 3) \text{ instead}$$

$$\text{ plus } \oplus \text{ on polynomials may operations on tables}$$

$$\text{ times } \odot$$

$$plus \oplus = coordinateurise addition$$

$$(S \oplus t)_{(\frac{k_{1}}{k_{2}}, \dots, \frac{k_{n}}{m})} = S_{(\frac{k_{1}}{k_{2}}, \dots, \frac{k_{n}}{m})} + t(\frac{k_{1}}{k_{2}}, \dots, \frac{k_{n}}{m})$$

$$adition in k$$

$$(S(\frac{k_{1}}{k_{2}}, \dots, \frac{k_{n}}{m}) \oplus \frac{1}{m})$$

$$o - dimensional linear space$$

$$S(\frac{k_{1}}{k_{2}}, \dots, \frac{k_{n}}{m}) \oplus \frac{1}{m}$$

$$n-times infinite sequence$$
Example :

$$f(x) = 2x^{3} + x + 3 \equiv (3 | 1 | 0 | 2 | 0 | 0, \dots, \infty)$$
  

$$g(x) = 4x^{4} + x^{3} + x^{2} \equiv (0 | 0 | 1 | 1 , 4 | 0, \dots, \infty)$$
  

$$(f \oplus g)(x) \equiv (3 | 1 | 1 | 3 | 4 | 0, \dots, \infty)$$

$$fines \ \odot = "shifting" of elements (algebraic ~ mon-linear operation)$$

$$(S \odot t)_{(k_{1}, k_{2}, \dots, k_{n})} = \sum_{\{(\alpha_{1}, \beta_{1}) \mid \alpha_{1}, \beta_{1} \in \mathbb{Z}_{20}^{n}, \alpha_{1} \neq \beta_{2} = (k_{1}, k_{2}, \dots, k_{n})\}$$

$$Example : f(x) = 2x^{3} + x + 3 \equiv (3 + 4 + 0 + 2 + 0 + 0 + 0 + 0 + 0 + 0)$$

$$g(x) = 4x^{4} + x^{3} + x^{2} \equiv (0 + 0 + 4 + 1 + 1 + 0 + 0 + 0)$$

$$(f \odot g)(x) = \sum_{\{(\alpha_{1}, 0)\}} f_{\alpha} + \sum_{\{(\alpha_{1}$$

Dual nature of polynomials - "or tables " & function  
Observation  
(i) 
$$f(x_{11}x_{21}...,x_n): k^n \longrightarrow k$$
 is a function  
(ii) Assume the field  $Z_z = (\xi \circ_1 13_1 \times +y \mod 2_1 \times \cdot y)$   
 $0+0=0 \mod 2=0 = \text{ neutrol element of the addition}$   
 $0+1=1 \mod 2=1$   
 $0+0=0 \ldots \text{ inverse of } 0 \text{ is } 0$   
 $1+1=2 \mod 2=0$   
 $1+1=0 \ldots \text{ inverse of } 1 \text{ is } 1$   
 $0\cdot0=0$   
 $0\cdot1=0$   
 $1\cdot1=1 = \text{ identity of the nultiplication, inverse of 1 is 1}$   
And now the punch line!  
 $x(x+1) = x^2 - x$  is NOT a zero polynomial, but  
 $0(0+1) = 0.1 = 0$  &  $1\cdot(1+1) = 1\cdot0 = 0$  is the zero function

## Excercises:

1. Is  $\mathbb{Z}_3$ , i.e. ({0,1,2},  $x + y \mod 3, x \cdot y$ ), a field?

2. Is there a non-zero polynomial in  $\mathbb{Z}_3$ , which is the zero function?