

# Advanced Robotics

## Lecture 5

## Algebraic equations & affine varieties

alg. equations

$$\underbrace{f_k(x_1, x_2, \dots, x_n)}_{\text{polynomial}} = 0 \quad k=1, \dots, s$$

monomial



$$f_k(x_1, x_2, \dots, x_n) = \sum_{\alpha} a_{\alpha} x^{\alpha}$$

$$x^{\alpha} = x_1^{\alpha_1} \cdot x_2^{\alpha_2} \cdots x_n^{\alpha_n} \quad \alpha = (\alpha_1, \alpha_2, \dots, \alpha_n) \quad \alpha_i \in \underbrace{\mathbb{Z}_{\geq 0}}$$

$a_{\alpha} \in k$  ... a field ( $\mathbb{R}, \mathbb{C}, \dots$ )

*nonnegative  
whole numbers*

$k[x_1, x_2, \dots, x_n]$  ≡ the set of all polynomials in  $x_1, x_2, \dots, x_n$

$(x_1, x_2, \dots, x_n) \in k^n$  ≡  $n$ -dimensional linear space

Example: variables  $x_1, x_2, x_3$

$$f = 2x_1^3 x_2^2 x_3 + \frac{3}{2} x_2^3 x_3^3 - 3x_1 x_2 x_3 + x_2^2$$

The convention:

$$f = 2x^{(3,2,1)} + \frac{3}{2}x^{(0,2,3)} - 3x^{(1,1,1)} + x^{(0,2,0)}$$

Diagram illustrating the components of a term:

- coefficient**: The numerical factor in front of the monomial.
- term**: The product of the coefficient and the monomial.
- monomial**: The variable part of the term, consisting of variables  $x_1, x_2, x_3$  with their respective exponents.
- exponent**: The power to which each variable is raised, represented by a tuple  $(\alpha_1, \alpha_2, \alpha_3)$ .

Total degree of  $f$ :  $\deg(f) = \max_{\alpha \in f} |\alpha|$ , where  $f = \sum_{\alpha} a_{\alpha} x^{\alpha}$ ,  $\alpha \in \mathbb{Z}_{\geq 0}^n$

$$|\alpha| = \alpha_1 + \alpha_2 + \dots + \alpha_n$$

Example:  $\deg(0) = -\infty$ , or undefined

$$\deg(f) = \max_{\alpha \in \{(3,2,1), (0,2,3), (0,2,0)\}} |\alpha| = |(3,2,1)| = 3+2+1 = 6$$

# Ring of polynomials

like  $\mathbb{R}$  without the division

plus  $\oplus$

times  $\odot$

$$\left\{ \begin{array}{l} k[x_1, x_2, \dots, x_n] \times k[x_1, x_2, \dots, x_n] \\ \text{Cartesian product} \end{array} \right\} \longrightarrow k[x_1, x_2, \dots, x_n]$$

i.e.  $f \oplus g = h$  or  $f \odot g = m$

Be  $f = \sum_{\alpha \in A} a_\alpha x^\alpha = \sum_{\alpha \in A \cup B} a'_\alpha x^\alpha$  where  $a'_\alpha = \begin{cases} a_\alpha, & \alpha \in A \\ 0, & \alpha \notin A \end{cases}$

$$g = \sum_{\beta \in B} b_\beta x^\beta = \sum_{\beta \in A \cup B} b'_\beta x^\beta$$

where  $b'_\beta = \begin{cases} b_\beta, & \beta \in B \\ 0, & \beta \notin B \end{cases}$

Define:

plus  $f \oplus g = \sum_{\gamma \in A \cup B} (a'_\gamma + b'_\gamma) x^\gamma$  ( $\sum_\alpha 0 \cdot x^\alpha$  is the ... zero polynomial)

times  $f \odot g = \sum_{\alpha \in A \cup B} \sum_{\beta \in A \cup B} (a'_\alpha \cdot b'_\beta) x^{(\alpha+\beta)}$  ( $1 \cdot x^{(0,0,\dots,0)} + 0 \cdot x^\alpha + \dots$  unit polynomial)

## Exercises:

1. Show that  $(k[x_1, \dots, k_n], \oplus, \odot)$  is a ring, i.e. that holds true

- (a)  $f \oplus g = g \oplus f$
- (b)  $(f \oplus g) \oplus h = f \oplus (g \oplus h)$
- (c)  $\exists 0 \in k[x_1, \dots, k_n]: f \oplus 0 = f$
- (d)  $\forall f \exists e \in k[x_1, \dots, k_n]: e \oplus f = 0$
- (e)  $f \odot g = g \odot f$
- (f)  $(f \odot g) \odot h = f \odot (g \odot h)$
- (g)  $\exists 1 \in k[x_1, \dots, k_n]: f \odot 1 = f$
- (h)  $f \odot (g \oplus h) = (f \odot g) \oplus (f \odot h)$   
for arbitrary  $f, g, h \in k[x_1, \dots, k_n]$ .

# Dual nature of polynomials - "∞ tables" vs functions

## Observations

- 1) there is **no limit** on the total degree of a polynomial
- 2) a polynomial has always a **finite** number of non-zero  $a_\alpha$ 's

$\Rightarrow$  polynomials  $\equiv$  **infinite** tables (sequences) with a **finite** number of non-zero entries

$$f = \sum_{\alpha \in A} a_\alpha x^\alpha = \underbrace{a_{(0,0,\dots,0)} x^{(0,0,\dots,0)}}_{\text{non-zero "coordinates" of } s(k_1, k_2, \dots, k_m), k_i \in \mathbb{Z}_{\geq 0}} + \dots + \underbrace{a_{(\alpha_1, \alpha_2, \dots, \alpha_m)} x^{(\alpha_1, \alpha_2, \dots, \alpha_m)}}_{+ \dots}$$

$$f \equiv s(\underbrace{k_1, k_2, \dots, k_m}_m\text{-times infinite sequence}) \quad s(k_1, k_2, \dots, k_m) = \begin{cases} a_\alpha & \text{if } \exists \alpha \in A : (k_1, k_2, \dots, k_m) = \alpha \\ 0 & \text{otherwise} \end{cases}$$

$k_i \in (1, 2, \dots, \infty) \equiv$  infinite sequence

Example:

$$f(x) = 2x^3 + x + 3 = 3x^0 + 1 \cdot x^1 + 0 \cdot x^2 + 2x^3 + 0 \cdot x^4 + 0 \cdot x^5 + \dots$$



$$\equiv (3, 1, 0, 2, 0, 0, \dots, \infty) \quad \dots \text{an infinite sequence}$$

$x^0 \ x^1 \ x^2 \ x^3 \ x^4 \ x^5 \ \dots$

with non-zero values  
on a finite number  
of coordinates

but we often write  $(2, 0, 1, 3)$  instead

plus  $\oplus$       on polynomials  $\rightsquigarrow$  operations on tables

times  $\odot$

plus  $\oplus \equiv$  coordinatewise addition

$$(s \oplus t)_{(k_1, k_2, \dots, k_n)} = s_{(k_1, k_2, \dots, k_n)} + t_{(k_1, k_2, \dots, k_n)}$$

addition in  $k$

$s_{(k_1, k_2, \dots, k_n)}$ :  $\oplus, \cdot$  element-wise multiplication by a scalar  $\in k$

$\underbrace{(s_{(k_1, k_2, \dots, k_n)})}_{\infty\text{-dimensional linear space}}$

$$s_{(k_1, k_2, \dots, k_n)} \in k^{(1, 2, \dots, \infty)^n}$$

$n$ -times infinite sequence

Example:

$$f(x) = 2x^3 + x + 3 \equiv (3, 1, 0, 2, 0, 0, \dots, \infty)$$

$$g(x) = 4x^4 + x^3 + x^2 \equiv (0, 0, 1, 1, 4, 0, \dots, \infty)$$

$$(f \oplus g)(x) \equiv (3, 1, 1, 3, 4, 0, \dots, \infty)$$

times  $\odot \equiv$  "shifting" of elements (algebraic  $\sim$  non-linear operation)

$$(S \odot t)_{(k_1, k_2, \dots, k_n)} = \sum_{\{(\alpha, \beta) \mid \alpha, \beta \in \mathbb{Z}_{\geq 0}^n, \alpha + \beta = (k_1, k_2, \dots, k_n)\}} s_\alpha \cdot t_\beta$$

Example:  $f(x) = 2x^3 + x + 3 \equiv (3, 1, 0, 2, 0, 0, \dots, \infty)$   
 $g(x) = 4x^4 + x^3 + x^2 \equiv (0, 0, 1, 1, 4, 0, 0, \dots, \infty)$

$$\begin{aligned}
 (f \odot g)(x) &\equiv \underbrace{\sum_{\{(0,0)\}} f_\alpha \cdot g_\beta}_{(f \odot g)(0)} + \underbrace{\sum_{\{(1,0), (0,1)\}} f_\alpha \cdot g_\beta}_{(f \odot g)(1)} + \underbrace{\sum_{\{(2,0), (1,1), (0,2)\}} f_\alpha \cdot g_\beta}_{(f \odot g)(2)} + \dots + \underbrace{\sum_{\{(7,0), (6,1), \dots, (0,7)\}} f_\alpha \cdot g_\beta}_{(f \odot g)(7)} \\
 &= (3 \cdot 0, 1 \cdot 0 + 3 \cdot 0, 0 \cdot 0 + 1 \cdot 0 + 3 \cdot 1, \dots, 0 \cdot 0 + 0 \cdot 0 + \dots + 3 \cdot 0, \dots) \\
 &= (0, 0, 3, 3+1, 12+1, 2+4, 2+8, 0, 0, \dots, \infty)
 \end{aligned}$$

Dual nature of polynomials - "∞ tables" & functions

Observation

(i)  $f(x_1, x_2, \dots, x_m) : k^n \rightarrow k$  is a function

(ii) Assume the field  $\mathbb{Z}_2 = (\{0, 1\}, x+y \bmod 2, x \cdot y)$

$0+0=0 \bmod 2=0$   $\equiv$  neutral element of the addition

$0+1=1 \bmod 2=1$   $0+0=0$  ... inverse of 0 is 0

$1+1=2 \bmod 2=0$   $1+1=0$  .... inverse of 1 is 1

$0 \cdot 0 = 0$

$0 \cdot 1 = 0$

$1 \cdot 1 = 1$   $\equiv$  identity of the multiplication, inverse of 1 is 1

And now the punch line!

$x(x+1) = x^2 - x$  is NOT a zero polynomial, but

$0 \cdot (0+1) = 0 \cdot 1 = 0$  &  $1 \cdot (1+1) = 1 \cdot 0 = 0$  is the zero function

## Excercises:

1. Is  $\mathbb{Z}_3$ , i.e.  $(\{0, 1, 2\}, x + y \bmod 3, x \cdot y)$ , a field?
2. Is there a non-zero polynomial in  $\mathbb{Z}_3$ , which is the zero function?

## Polynomials in one variable

leading term: a non-zero polynomial

$$f(x) = a_m x^m + a_{m-1} x^{m-1} + \dots + a_0 \in k[x]$$

$a_m \neq 0$

$$\text{LT}(f) = a_m x^m \equiv \text{the leading term}$$

Example:

$$f = 2x^3 - 4x + 3 \Rightarrow \text{LT}(f) = 2x^3$$

## Division of terms

$\alpha, \beta \in \mathbb{Z}_{\geq 0}^n$ ,  $a_\alpha, b_\beta \in k$ ,  $x^\alpha, x^\beta \in k[x_1, \dots, x_n]$  monomials

$a_\alpha x^\alpha$  divides  $b_\beta x^\beta$   $\stackrel{\text{def}}{=}$   $\beta_i - \alpha_i \geq 0$ ,  $i = 1, \dots, n$

If  $a_\alpha x^\alpha$  divides  $b_\beta x^\beta$ , then there is exactly one monomial

$$c_\gamma x^\gamma = \frac{b_\beta}{a_\alpha} \cdot x^{\beta - \alpha}$$

such that  $b_\beta x^\beta = a_\alpha x^\alpha \cdot c_\gamma x^\gamma$

"Division" of polynomials in one variable

polynomials **cannot** be divided but can be "divided"

$$f : g \stackrel{\text{def}}{\equiv} f = qg + r, \quad r=0 \vee \deg(r) < \deg(g)$$

Example  $f = 2x^3 - 4x + 3, \quad g(x) = x - 1$

$$\begin{aligned} f : g &\equiv 2x^3 - 4x + 3 = 2x^2(x-1) + 2x^2 - 4x + 3 = \\ &= (2x^2 + 2x)(x-1) - 2x + 3 = \underbrace{(2x^2 + 2x - 2)(x-1)}_q + \underbrace{1}_r \end{aligned}$$

Notice that:  $\deg(f) = \deg(\text{LT}(f))$

$\text{LT}(g)$  divides  $\text{LT}(f) \Leftrightarrow \deg(\text{LT}(g)) \leq \deg(\text{LT}(f)) \Leftrightarrow \deg(g) \leq \deg(f)$

$\text{LT}(g)$  divides  $\text{LT}(f) \Leftrightarrow \deg(g) \leq \deg(f)$

## "Division Theorem"

Let  $k$  be a field and  $g$  be a non-zero polynomial in  $k[x]$ .

- (i) Then every  $f \in k[x]$  can be written as

$$f = qg + r$$

where  $q, r \in k[x]$ , and either

$$r=0 \text{ or } \deg(r) < \deg(g)$$

- (ii) Furthermore,  $q$  and  $r$  are unique.