

# **CAN Device Driver Internals**

**Pavel Pisa**

**pisa@cmp.felk.cvut.cz**

## **CAN Device Driver Internals**

by Pavel Pisa

Copyright © 2002 by Pavel Pisa

The LinCAN is an implementation of the Linux device driver supporting more CAN controller chips and many CAN interface boards. Its implementation has long history already. The OCERA version of the driver adds new features, continuous enhancements and reimplementations of structure of the driver. The usage of the driver is tightly coupled to the virtual CAN API interface component which hides driver low level interface to the application programmers.

This book describes internals of the LinCAN driver.

This documentation is free software; you can redistribute it and/or modify it under the terms of the GNU General Public License as published by the Free Software Foundation; either version 2 of the License, or (at your option) any later version.

This program is distributed in the hope that it will be useful, but WITHOUT ANY WARRANTY; without even the implied warranty of MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the GNU General Public License for more details.

You should have received a copy of the GNU General Public License along with this program; if not, write to the Free Software Foundation, Inc., 59 Temple Place, Suite 330, Boston, MA 02111-1307 USA

For more details see the file COPYING in the source distribution of LinCAN driver.

# Table of Contents

1. Introduction .....	1
2. LinCAN Driver API.....	2
Driver API Overview.....	2
CAN Driver File Operations.....	3
open .....	3
close .....	3
read .....	4
write.....	4
3. Driver Internal Documentation.....	6
Basic Driver Data Structures.....	6
struct canhardware_t.....	6
struct candevice_t .....	6
struct chip_t .....	8
struct msgobj_t.....	9
struct canuser_t .....	10
struct hwspecops_t .....	11
struct chipspecops_t .....	12
Board Support Functions .....	14
template_request_io .....	14
template_release_io .....	15
template_reset .....	15
template_init_hw_data .....	16
template_init_chip_data .....	17
template_init_obj_data.....	18
template_program_irq .....	18
template_write_register .....	19
template_read_register .....	20
Chip Support Functions .....	20
sja1000p_enable_configuration.....	21
sja1000p_disable_configuration.....	21
sja1000p_chip_config .....	21
sja1000p_extended_mask .....	22
sja1000p_baud_rate .....	23
sja1000p_read .....	23
sja1000p_pre_read_config .....	24
sja1000p_pre_write_config .....	25
sja1000p_send_msg .....	25
sja1000p_check_tx_stat .....	26
sja1000p_set_btregs .....	27
sja1000p_start_chip.....	27
sja1000p_stop_chip .....	28
sja1000p_remote_request.....	29
sja1000p_standard_mask .....	29
sja1000p_clear_objects.....	30
sja1000p_config_irqs.....	30
sja1000p_irq_write_handler .....	31
sja1000p_irq_handler .....	32
sja1000p_wakeup_tx .....	33
CAN Queues Structures and Functions.....	33
struct canque_slot_t.....	33
struct canque_fifo_t .....	34
canque_fifo_get_inslot .....	35
canque_fifo_put_inslot .....	35
canque_fifo_abort_inslot .....	36
canque_fifo_test_outslot .....	36
canque_fifo_free_outslot .....	37
canque_fifo_again_outslot .....	37
struct canque_edge_t.....	38
struct canque_ends_t.....	39
canque_notify_inends .....	40
canque_notify_outends .....	41
canque_notify_bothends .....	41

canque_activate_edge .....	42
canque_filtid2internal .....	42
canque_fifo_flush_slots.....	43
canque_fifo_init_slots .....	43
canque_fifo_done.....	44
canque_get_inslot .....	44
canque_get_inslot4id.....	45
canque_put_inslot .....	46
canque_abort_inslot .....	47
canque_filter_msg2edges .....	47
canque_test_outslot .....	48
canque_free_outslot.....	49
canque_again_outslot .....	49
canque_set_filt.....	50
canque_flush.....	51
canqueue_ends_init_gen .....	51
canqueue_notify_kern.....	52
canqueue_ends_init_kern.....	52
canque_get_inslot4id_wait_kern.....	53
canque_get_outslot_wait_kern.....	53
canque_sync_wait_kern.....	54
canque_new_edge_kern .....	55
canqueue_connect_edge .....	55
canqueue_disconnect_edge .....	56
canqueue_disconnect_edge_kern.....	56
canqueue_ends_done_kern.....	57

## **Chapter 1. Introduction**

The LinCAN driver is the loadable module for the Linux kernel which implements CAN driver. The driver communicates and controls one or more CAN controllers chips. The each chip/CAN interface is represented to the applications as one or more CAN message objects through the character device interface. The application can open the character device and use read/write system calls for CAN messages transmission or reception through the connected message object. The parameters of the message object can be modified by the IOCTL system call. The closing of the character device releases resources allocated by the application.

The present version of the driver supports three most common CAN controllers:

- Intel i82527 chips
- Philips 82c200 chips
- Philips SJA1000 chips in standard and PeliCAN mode

The intelligent CAN/CANopen cards should be supported by future versions. One of such cards is P-CAN series of cards produced by Unicontrols. The driver contains support for more than ten CAN cards basic types with different combinations of the above mentioned chips. Not all card types are held by OCERA members, but CTU has and tested more SJA1000 type cards and will test some i82527 cards in near future.

## Chapter 2. LinCAN Driver API

### Driver API Overview

Each driver is a subsystem which has no direct application level API. The operating system is responsible for user space calls transformation into driver functions calls or dispatch routines invocations. The CAN driver is implemented as a character device with the standard device node names `/dev/can0`, `/dev/can1`, etc. The application program communicates with the driver through the standard system low level input/output primitives (`open`, `close`, `read`, `write`, `select` and `ioctl`). The CAN driver convention of usage of these functions is described in the next subsection.

The `read` and `write` functions need to transfer one or more CAN messages. The structure `canmsg_t` is defined for this purpose and is defined in include file `can/can.h`. The `canmsg_t` structure has next fields:

```
struct canmsg_t {
    short flags;
    int cob;
    unsigned long id;
    unsigned long timestamp;
    unsigned int length;
    unsigned char
    data[CAN_MSG_LENGTH];
} PACKED;
```

#### flags

The flags field holds information about message type. The bit `MSG_RTR` marks remote transmission request messages. Writing of such message into the CAN message object handle results in transmission of the RTR message. The RTR message can be received by the `read` call if no buffer with corresponding ID is pre-filled in the driver. The bit `MSG_EXT` indicates that the message with extended (bit 29 set) ID will be send or was received. The bit `MSG_OVR` is intended for fast indication of the reception message queue overfill.

#### cob

The field reserved for a holding message communication object number. It could be used for serialization of received messages from more message object into one message queue in the future.

#### id

CAN message ID.

#### timestamp

The field intended for storing of the message reception time.

#### length

The number of the data bytes send or received in the CAN message. The number of data load bytes is from 0 to 8.

#### data

The byte array holding message data.

As was mentioned above, direct communication with the driver through system calls is not encouraged because this interface is partially system dependent and cannot be ported to all environments. The suggested alternative is to use OCERA provided VCA library which defines the portable and clean interface to the CAN driver implementation.

The other issue is addition of the support for new CAN interface boards and CAN controller chips. The subsection Board Support Functions describes template functions, which needs to be implemented for newly supported board. The template of board support can be found in the file `src/template.c`.

The other task for more brave souls is addition of the support for the unsupported chip type. The source supporting the SJA1000 chip in the PeliCAN mode can serve as an example. The full source of this chip support is stored in the file `src/sja1000p.c`. The subsection Chip Support Functions describes basic functions necessary for the new chip support.

## CAN Driver File Operations

### open

#### Name

`open` — message communication object open system call

#### Synopsis

```
int open (const char * pathname, int flags);
```

#### Arguments

*pathname*

The path to driver device node is specified there. The conventional device names for Linux CAN driver are `/dev/can0`, `/dev/can1`, etc.

*flags*

flags modifying style of open call. The standard `O_RDWR` mode should be used for CAN device. The mode `O_NOBLOCK` can be used with driver as well. This mode results in immediate return of read and write.

#### Description

Returns negative number in the case of error. Returns the file descriptor for named CAN message object in other cases.

### close

#### Name

`close` — message communication object close system call

#### Synopsis

```
int close (int fd);
```

## Arguments

## Chapter 2. LinCAN Driver API

*fd*

file descriptor to opened can message communication object

## Description

Returns negative number in the case of error.

## read

### Name

read — reads received CAN messages from message object

### Synopsis

```
ssize_t read(int fd, void * buf, size_t count);
```

## Arguments

*fd*

file descriptor to opened can message communication object

*buf*

pointer to array of canmsg\_t structures.

*count*

size of message array buffer in number of bytes

## Description

Returns negative value in the case of error else returns number of read bytes which is multiple of canmsg\_t structure size.

## write

### Name

write — writes CAN messages to message object for transmission

```
ssize_t write(int fd, const void * buf, size_t count);
```

## Arguments

*fd*

file descriptor to opened can message communication object

*buf*

pointer to array of canmsg\_t structures.

*count*

size of message array buffer in number of bytes. The parameter informs driver about number of messages prepared for transmission and should be multiple of canmsg\_t structure size.

## Description

Returns negative value in the case of error else returns number of bytes successfully stored into message object transmission queue. The positive returned number is multiple of canmsg\_t structure size.

## Chapter 3. Driver Internal Documentation

### Basic Driver Data Structures

#### **struct canhardware\_t**

##### **Name**

`struct canhardware_t` — structure representing pointers to all CAN boards

##### **Synopsis**

```
struct canhardware_t {
    int nr_boards;
    struct rtr_id * rtr_queue;
    spinlock_t rtr_lock;
    struct candevice_t ** candevice;
};
```

##### **Members**

`nr_boards`

number of present boards

`rtr_queue`

RTR - remote transmission request queue (expect some changes there)

`rtr_lock`

locking for RTR queue

`candevice`

array of pointers to CAN devices/boards

#### **struct candevice\_t**

##### **Name**

`struct candevice_t` — CAN device/board structure

##### **Synopsis**

```
struct candevice_t {
    char * hwname;
    int candev_idx;
    unsigned long io_addr;
    unsigned long res_addr;
    unsigned long dev_base_addr;
    unsigned int flags;
    int nr_all_chips;
    int nr_82527_chips;
    int nr_sja1000_chips;
    struct chip_t ** chip;
    struct hwspecops_t * hwspecops;
```

```
    struct canhardware_t * hosthardware; Chapter 3. Driver Internal Documentation
};
```

## Members

hwname  
text string with board type

candev\_idx  
board index in canhardware\_t.candevice[]

io\_addr  
IO/physical MEM address

res\_addr  
optional reset register port

dev\_base\_addr  
CPU translated IO/virtual MEM address

flags  
board flags: PROGRAMMABLE\_IRQ .. interrupt number can be programmed into board

nr\_all\_chips  
number of chips present on the board

nr\_82527\_chips  
number of Intel 8257 chips

nr\_sja1000\_chips  
number of Philips SJA100 chips

chip  
array of pointers to the chip structures

hwspecops  
pointer to board specific operations

hosthardware\_p  
pointer to the root hardware structure

## Description

The structure represent configuration and state of associated board. The driver infrastructure prepares this structure and calls board type specific `board_register` function. The board support provided register function fills right function pointers in `hwspecops` structure. Then driver setup calls functions `init_hw_data`, `init_chip_data`, `init_obj_data` and `program_irq`. Function `init_hw_data` and `init_chip_data` have to specify number and types of connected chips or objects respectively. The use of `nr_all_chips` is preferred over use of fields `nr_82527_chips` and `nr_sja1000_chips` in the board non-specific functions. The `io_addr` and `dev_base_addr` is filled from module parameters to the same value. The `request_io` function can fix-up `dev_base_addr` field if virtual address is different than bus address.

## Name

struct chip\_t — CAN chip state and type information

## Synopsis

```
struct chip_t {
    char * chip_type;
    int chip_idx;
    int chip_irq;
    unsigned long chip_base_addr;
    unsigned int flags;
    int clock;
    void (* write_register (unsigned char data,unsigned long address));
    unsigned (* read_register (unsigned long address));
    unsigned short sja_cdr_reg;
    unsigned short sja_ocr_reg;
    unsigned short int_cpu_reg;
    unsigned short int_clk_reg;
    unsigned short int_bus_reg;
    struct msgobj_t ** msgobj;
    struct chipspecops_t * chipspecops;
    struct candevice_t * hostdevice;
    int max_objects;
};
```

## Members

chip\_type

text string describing chip type

chip\_idx

index of the chip in candevice\_t.chip[] array

chip\_irq

chip interrupt number if any

chip\_base\_addr

chip base address in the CPU IO or virtual memory space

flags

chip flags: CHIP\_CONFIGURED .. chip is configured, CHIP\_SEGMENTED .. access to the chip is segmented (mainly for i82527 chips)

clock

chip base clock frequency in Hz

write\_register

write chip register function copy -

read\_register

read chip register function copy

sja\_cdr\_reg

SJA specific register - holds hardware specific options for the Clock Divider register. Options defined in the sja1000.h file: CDR\_CLKOUT\_MASK, CDR\_CLK\_OFF, CDR\_RXINPEN, CDR\_CBP, CDR\_PELICAN

SJA specific register - hold hardware specific options for the Output Control register. Options defined in the sja1000.h file: OCR\_MODE\_BIPHASE, OCR\_MODE\_TEST, OCR\_MODE\_NORMAL, OCR\_MODE\_CLOCK, OCR\_TX0\_LH, OCR\_TX1\_ZZ.

#### int\_cpu\_reg

Intel specific register - holds hardware specific options for the CPU Interface register. Options defined in the i82527.h file: iCPU\_CEN, iCPU\_MUX, iCPU\_SLP, iCPU\_PWD, iCPU\_DMC, iCPU\_DSC, iCPU\_RST.

#### int\_clk\_reg

Intel specific register - holds hardware specific options for the Clock Out register. Options defined in the i82527.h file: iCLK\_CD0, iCLK\_CD1, iCLK\_CD2, iCLK\_CD3, iCLK\_SL0, iCLK\_SL1.

#### int\_bus\_reg

Intel specific register - holds hardware specific options for the Bus Configuration register. Options defined in the i82527.h file: iBUS\_DR0, iBUS\_DR1, iBUS\_DT1, iBUS\_POL, iBUS\_CBY.

#### msgobj

array of pointers to individual communication objects

#### chipspecops

pointer to the set of chip specific object filled by init\_chip\_data function

#### hostdevice

pointer to chip hosting board

#### max\_objects

maximal number of communication objects connected to this chip

## Description

The fields `write_register` and `read_register` are copied from corresponding fields from `hwspecops` structure (`chip->hostdevice->hwspecops->write_register` and `chip->hostdevice->hwspecops->read_register`) to speedup `can_write_reg` and `can_read_reg` functions.

## struct msgobj\_t

### Name

`struct msgobj_t` — structure holding communication object state

### Synopsis

```
struct msgobj_t {
    unsigned long obj_base_addr;
    unsigned int minor;
    unsigned int object;
    unsigned long flags;
    int ret;
    struct canque_ends_t * qends;
    struct canque_edge_t * tx_qedge;
    struct canque_slot_t * tx_slot;
    int tx_retry_cnt;
```

```

    struct canmsg_t rx_msg;
    struct chip_t * hostchip;
    atomic_t obj_used;
    struct list_head obj_users;
} ;

```

## Members

obj\_base\_addr

minor

associated device minor number

object

object number in chip\_t structure +1

flags

message object flags

ret

field holding status of the last Tx operation

qends

pointer to message object corresponding ends structure

tx\_qedge

edge corresponding to transmitted message

tx\_slot

slot holding transmitted message, slot is taken from canque\_test\_outslot call and is freed by canque\_free\_outslot or rescheduled canque\_again\_outslot

tx\_retry\_cnt

transmission attempt counter

rx\_msg

temporary storage to hold received messages before calling to canque\_filter\_msg2edges

hostchip

pointer to the &chip\_t structure this object belongs to

obj\_used

counter of users (associated file structures for Linux userspace clients) of this object

obj\_users

list of user structures of type &canuser\_t.

**Name**

struct canuser\_t — structure holding CAN user/client state

**Synopsis**

```
struct canuser_t {
    struct list_head peers;
    struct canque_ends_t * qends;
    struct file * file;
    struct msgobj_t * msgobj;
    struct canque_edge_t * rx_edge0;
    int magic;
};
```

**Members**

peers

for connection into list of object users

qends

pointer to the ends structure corresponding for this user

file

pointer to open device file state structure

msgobj

communication object the user is connected to

rx\_edge0

default receive queue for filter IOCTL

magic

magic number to check consistency when pointer is retrieved from file private field

**struct hwspecops\_t****Name**

struct hwspecops\_t — hardware/board specific operations

**Synopsis**

```
struct hwspecops_t {
    int (* request_io (struct candevice_t *candev);
    int (* release_io (struct candevice_t *candev);
    int (* reset (struct candevice_t *candev);
    int (* init_hw_data (struct candevice_t *candev);
    int (* init_chip_data (struct candevice_t *candev, int chipnr);
    int (* init_obj_data (struct chip_t *chip, int objnr);
    int (* program_irq (struct candevice_t *candev);
    void (* write_register (unsigned char data,unsigned long address);
```

Chapter 3, Driver Internal Documentation

```
    unsigned (* read_register (unsigned long address));
```

## Members

request\_io  
    reserve io or memory range for can board

release\_io  
    free reserved io memory range

reset  
    hardware reset routine

init\_hw\_data  
    called to initialize &candev\_t structure, mainly res\_addr, nr\_all\_chips,  
    nr\_i82527\_chips, nr\_sja1000\_chips and flags fields

init\_chip\_data  
    called initialize each &chip\_t structure, mainly chip\_type,  
    chip\_base\_addr, clock and chip specific registers. It is responsible to setup  
    &chip\_t->chipspecops functions for non-standard chip types (type other  
    than "i82527", "sja1000" or "sja1000p")

init\_obj\_data  
    called initialize each &msgobj\_t structure, mainly obj\_base\_addr field.

program\_irq  
    program interrupt generation hardware of the board if flag PROGRAMMABLE\_IRQ  
    is present for specified device/board

write\_register  
    low level write register routine

read\_register  
    low level read register routine

## struct chipspecops\_t

### Name

struct chipspecops\_t — can controller chip specific operations

### Synopsis

```
struct chipspecops_t {  
    int (* chip_config (struct chip_t *chip));  
    int (* baud_rate (struct chip_t *chip, int rate, int clock, int sjw,int sampl_pt, int);  
    int (* standard_mask (struct chip_t *chip, unsigned short code,unsigned short mask);  
    int (* extended_mask (struct chip_t *chip, unsigned long code,unsigned long mask);  
    int (* message15_mask (struct chip_t *chip, unsigned long code,unsigned long mask);  
    int (* clear_objects (struct chip_t *chip);  
    int (* config_irqs (struct chip_t *chip, short irqs);  
    int (* pre_read_config (struct chip_t *chip, struct msgobj_t *obj);  
    int (* pre_write_config (struct chip_t *chip, struct msgobj_t *obj,struct canmsg_t *m);
```

```

int (* send_msg (struct chip_t *chip, struct msgobj_t *obj, struct canmsg_t *msg);
int (* remote_request (struct chip_t *chip, struct msgobj_t *obj));
int (* check_tx_stat (struct chip_t *chip));
int (* wakeup_tx (struct chip_t *chip, struct msgobj_t *obj));
int (* enable_configuration (struct chip_t *chip));
int (* disable_configuration (struct chip_t *chip));
int (* set_btregs (struct chip_t *chip, unsigned short btr0,unsigned short btr1);
int (* start_chip (struct chip_t *chip);
int (* stop_chip (struct chip_t *chip);
irqreturn_t (* irq_handler (int irq, void *dev_id, struct pt_regs *regs);
};
```

### *Chapter 3. Driver Internal Documentation*

## Members

chip\_config

CAN chip configuration

baud\_rate

set communication parameters

standard\_mask

setup of mask for message filtering

extended\_mask

setup of extended mask for message filtering

message15\_mask

set mask of i82527 message object 15

clear\_objects

clears state of all message object residing in chip

config\_irqs

tunes chip hardware interrupt delivery

pre\_read\_config

prepares message object for message reception

pre\_write\_config

prepares message object for message transmission

send\_msg

initiate message transmission

remote\_request

configures message object and asks for RTR message

check\_tx\_stat

checks state of transmission engine

wakeup\_tx

wakeup TX processing

enable\_configuration

enable chip configuration mode

disable\_configuration

disable chip configuration mode

set_btregs	<i>Chapter 3. Driver Internal Documentation</i>
configures bitrate registers	
start_chip	
starts chip message processing	
stop_chip	
stops chip message processing	
irq_handler	
interrupt service routine	

## Board Support Functions

The functions, which should be implemented for each supported board, are described in the next section. The functions are prefixed by boardname. The prefix *template* has been selected for next description.

### **template\_request\_io**

#### **Name**

`template_request_io` — reserve io or memory range for can board

#### **Synopsis**

```
int template_request_io (struct candevice_t * candevo);
```

#### **Arguments**

*candevo*

pointer to candevice/board which asks for io. Field *io\_addr* of *candevo* is used in most cases to define start of the range

#### **Description**

The function `template_request_io` is used to reserve the io-memory. If your hardware uses a dedicated memory range as hardware control registers you will have to add the code to reserve this memory as well. `IO_RANGE` is the io-memory range that gets reserved, please adjust according your hardware. Example: `#define IO_RANGE 0x100` for i82527 chips or `#define IO_RANGE 0x20` for sja1000 chips in basic CAN mode.

#### **Return Value**

The function returns zero on success or `-ENODEV` on failure

## template\_release\_io

### Name

template\_release\_io — free reserved io memory range

### Synopsis

```
int template_release_io (struct candevice_t * candevo);
```

### Arguments

*candevo*

pointer to candevice/board which releases io

### Description

The function template\_release\_io is used to free reserved io-memory. In case you have reserved more io memory, don't forget to free it here. IO\_RANGE is the io-memory range that gets released, please adjust according your hardware. Example:  
#define IO\_RANGE 0x100 for i82527 chips or #define IO\_RANGE 0x20 for sja1000 chips in basic CAN mode.

### Return Value

The function always returns zero

**File**

## template\_reset

### Name

template\_reset — hardware reset routine

## Synopsis

## Chapter 3. Driver Internal Documentation

```
int template_reset (struct candevice_t * cande);
```

## Arguments

*cande*

Pointer to candevice/board structure

## Description

The function `template_reset` is used to give a hardware reset. This is rather hardware specific so I haven't included example code. Don't forget to check the reset status of the chip before returning.

## Return Value

The function returns zero on success or `-ENODEV` on failure

## File

`src/template.c`

## **template\_init\_hw\_data**

### Name

`template_init_hw_data` — Initialize hardware cards

## Synopsis

```
int template_init_hw_data (struct candevice_t * cande);
```

## Arguments

*cande*

Pointer to candevice/board structure

## Description

The function `template_init_hw_data` is used to initialize the hardware structure containing information about the installed CAN-board. `RESET_ADDR` represents the io-address of the hardware reset register. `NR_82527` represents the number of Intel 82527 chips on the board. `NR_SJA1000` represents the number of Philips sj1000 chips on the board. The `flags` entry can currently only be `PROGRAMMABLE_IRQ` to indicate that the hardware uses programmable interrupts.

## Return Value

## Chapter 3. Driver Internal Documentation

The function always returns zero

## File

src/template.c

# template\_init\_chip\_data

## Name

template\_init\_chip\_data — Initialize chips

## Synopsis

```
int template_init_chip_data (struct candevice_t * candeve, int chipnr);
```

## Arguments

*candeve*

Pointer to candevice/board structure

*chipnr*

Number of the CAN chip on the hardware card

## Description

The function `template_init_chip_data` is used to initialize the hardware structure containing information about the CAN chips. `CHIP_TYPE` represents the type of CAN chip. `CHIP_TYPE` can be “i82527” or “sja1000”. The `chip_base_addr` entry represents the start of the ‘official’ memory map of the installed chip. It’s likely that this is the same as the `io_addr` argument supplied at module loading time. The `clock` entry holds the chip clock value in Hz. The entry `sja_cdr_reg` holds hardware specific options for the Clock Divider register. Options defined in the `sja1000.h` file: `CDR_CLKOUT_MASK`, `CDR_CLK_OFF`, `CDR_RXINPEN`, `CDR_CBP`, `CDR_PELICAN`. The entry `sja_ocr_reg` holds hardware specific options for the Output Control register. Options defined in the `sja1000.h` file: `OCR_MODE_BIPHASE`, `OCR_MODE_TEST`, `OCR_MODE_NORMAL`, `OCR_MODE_CLOCK`, `OCR_TX0_LH`, `OCR_TX1_ZZ`. The entry `int_clk_reg` holds hardware specific options for the Clock Out register. Options defined in the `i82527.h` file: `iCLK_CD0`, `iCLK_CD1`, `iCLK_CD2`, `iCLK_CD3`, `iCLK_SL0`, `iCLK_SL1`. The entry `int_bus_reg` holds hardware specific options for the Bus Configuration register. Options defined in the `i82527.h` file: `iBUS_DR0`, `iBUS_DR1`, `iBUS_DT1`, `iBUS_POL`, `iBUS_CBY`. The entry `int_cpu_reg` holds hardware specific options for the cpu interface register. Options defined in the `i82527.h` file: `iCPU_CEN`, `iCPU_MUX`, `iCPU_SLP`, `iCPU_PWD`, `iCPU_DMC`, `iCPU_DSC`, `iCPU_RST`.

## Return Value

## Chapter 3. Driver Internal Documentation

The function always returns zero

## File

src/template.c

# template\_init\_obj\_data

## Name

template\_init\_obj\_data — Initialize message buffers

## Synopsis

```
int template_init_obj_data (struct chip_t * chip, int objnr);
```

## Arguments

*chip*

Pointer to chip specific structure

*objnr*

Number of the message buffer

## Description

The function `template_init_obj_data` is used to initialize the hardware structure containing information about the different message objects on the CAN chip. In case of the sja1000 there's only one message object but on the i82527 chip there are 15. The code below is for a i82527 chip and initializes the object base addresses. The entry `obj_base_addr` represents the first memory address of the message object. In case of the sja1000 `obj_base_addr` is taken the same as the chips base address. Unless the hardware uses a segmented memory map, flags can be set zero.

## Return Value

The function always returns zero

## File

src/template.c

**Name**

`template_program_irq` — program interrupts

**Synopsis**

```
int template_program_irq (struct candevice_t * candevice);
```

**Arguments**

*candevice*

Pointer to candevice/board structure

**Description**

The function `template_program_irq` is used for hardware that uses programmable interrupts. If your hardware doesn't use programmable interrupts you should not set the `candevices_t->flags` entry to `PROGRAMMABLE_IRQ` and leave this function unedited. Again this function is hardware specific so there's no example code.

**Return value**

The function returns zero on success or `-ENODEV` on failure

**File**

`src/template.c`

**template\_write\_register****Name**

`template_write_register` — Low level write register routine

**Synopsis**

```
void template_write_register (unsigned char data, unsigned long address);
```

**Arguments**

*data*

data to be written

*address*

*Chapter 3. Driver Internal Documentation*

memory address to write to

## Description

The function `template_write_register` is used to write to hardware registers on the CAN chip. You should only have to edit this function if your hardware uses some specific write process.

## Return Value

The function does not return a value

## File

`src/template.c`

# **template\_read\_register**

## Name

`template_read_register` — Low level read register routine

## Synopsis

```
unsigned template_read_register (unsigned long address);
```

## Arguments

*address*

memory address to read from

## Description

The function `template_read_register` is used to read from hardware registers on the CAN chip. You should only have to edit this function if your hardware uses some specific read process.

## Return Value

The function returns the value stored in `address`

## File

`src/template.c`

The controller chip specific functions are described in the next section. The functions should be prefixed by chip type. Because documentation of chip functions has been retrieved from the actual SJA1000 PeliCAN support, the function prefix is *sja1000p*.

### **sja1000p\_enable\_configuration**

#### **Name**

`sja1000p_enable_configuration` — enable chip configuration mode

#### **Synopsis**

```
int sja1000p_enable_configuration (struct chip_t * chip);
```

#### **Arguments**

*chip*

pointer to chip state structure

### **sja1000p\_disable\_configuration**

#### **Name**

`sja1000p_disable_configuration` — disable chip configuration mode

#### **Synopsis**

```
int sja1000p_disable_configuration (struct chip_t * chip);
```

#### **Arguments**

*chip*

pointer to chip state structure

### **sja1000p\_chip\_config**

#### **Name**

`sja1000p_chip_config` — can chip configuration

## Synopsis

## Chapter 3. Driver Internal Documentation

```
int sja1000p_chip_config (struct chip_t * chip);
```

## Arguments

*chip*

pointer to chip state structure

## Description

This function configures chip and prepares it for message transmission and reception. The function resets chip, resets mask for acceptance of all messages by call to `sja1000p_extended_mask` function and then computes and sets baudrate with use of function `sja1000p_baud_rate`.

## Return Value

negative value reports error.

## File

src/sja1000p.c

# **sja1000p\_extended\_mask**

## Name

`sja1000p_extended_mask` — setup of extended mask for message filtering

## Synopsis

```
int sja1000p_extended_mask (struct chip_t * chip, unsigned long code,  
                           unsigned long mask);
```

## Arguments

*chip*

pointer to chip state structure

*code*

can message acceptance code

*mask*

can message acceptance mask

## Return Value

Chapter 3. Driver Internal Documentation

negative value reports error.

## File

src/sja1000p.c

# sja1000p\_baud\_rate

## Name

sja1000p\_baud\_rate — set communication parameters.

## Synopsis

```
int sja1000p_baud_rate (struct chip_t * chip, int rate, int clock, int  
                        sjw, int samp1_pt, int flags);
```

## Arguments

*chip*

pointer to chip state structure

*rate*

baud rate in Hz

*clock*

frequency of sja1000 clock in Hz (ISA osc is 14318000)

*sjw*

synchronization jump width (0-3) prescaled clock cycles

*samp1\_pt*

sample point in % (0-100) sets (TSEG1+1)/(TSEG1+TSEG2+2) ratio

*flags*

fields BTR1\_SAM, OCMODE, OCPL, OCTP, OCTN, CLK\_OFF, CBP

## Return Value

negative value reports error.

## File

src/sja1000p.c

**Name**

`sja1000p_read` — reads and distributes one or more received messages

**Synopsis**

```
void sja1000p_read (struct chip_t * chip, struct msgobj_t * obj);
```

**Arguments**

*chip*

pointer to chip state structure

*obj*

pointer to CAN message queue information

**File**

src/sja1000p.c

**sja1000p\_pre\_read\_config****Name**

`sja1000p_pre_read_config` — prepares message object for message reception

**Synopsis**

```
int sja1000p_pre_read_config (struct chip_t * chip, struct msgobj_t *  
    obj);
```

**Arguments**

*chip*

pointer to chip state structure

*obj*

pointer to message object state structure

## Return Value

## Chapter 3. Driver Internal Documentation

negative value reports error. Positive value indicates immediate reception of message.

## File

src/sja1000p.c

# sja1000p\_pre\_write\_config

## Name

`sja1000p_pre_write_config` — prepares message object for message transmission

## Synopsis

```
int sja1000p_pre_write_config (struct chip_t * chip, struct msgobj_t *  
obj, struct canmsg_t * msg);
```

## Arguments

*chip*

pointer to chip state structure

*obj*

pointer to message object state structure

*msg*

pointer to CAN message

## Description

This function prepares selected message object for future initiation of message transmission by `sja1000p_send_msg` function. The CAN message data and message ID are transferred from *msg* slot into chip buffer in this function.

## Return Value

negative value reports error.

## File

src/sja1000p.c

**Name**

`sja1000p_send_msg` — initiate message transmission

**Synopsis**

```
int sja1000p_send_msg (struct chip_t * chip, struct msgobj_t * obj,  
struct canmsg_t * msg);
```

**Arguments**

*chip*

pointer to chip state structure

*obj*

pointer to message object state structure

*msg*

pointer to CAN message

**Description**

This function is called after `sja1000p_pre_write_config` function, which prepares data in chip buffer.

**Return Value**

negative value reports error.

**File**

`src/sja1000p.c`

**sja1000p\_check\_tx\_stat****Name**

`sja1000p_check_tx_stat` — checks state of transmission engine

**Synopsis**

```
int sja1000p_check_tx_stat (struct chip_t * chip);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*chip*

pointer to chip state structure

## Return Value

negative value reports error. Positive return value indicates transmission under way status. Zero value indicates finishing of all issued transmission requests.

## File

src/sja1000p.c

# sja1000p\_set\_btregs

## Name

sja1000p\_set\_btregs — configures bitrate registers

## Synopsis

```
int sja1000p_set_btregs (struct chip_t * chip, unsigned short btr0,  
unsigned short btr1);
```

## Arguments

*chip*

pointer to chip state structure

*btr0*

bitrate register 0

*btr1*

bitrate register 1

## Return Value

negative value reports error.

## File

src/sja1000p.c

**Name**

`sja1000p_start_chip` — starts chip message processing

**Synopsis**

```
int sja1000p_start_chip (struct chip_t * chip);
```

**Arguments**

*chip*

pointer to chip state structure

**Return Value**

negative value reports error.

**File**

`src/sja1000p.c`

**sja1000p\_stop\_chip****Name**

`sja1000p_stop_chip` — stops chip message processing

**Synopsis**

```
int sja1000p_stop_chip (struct chip_t * chip);
```

**Arguments**

*chip*

pointer to chip state structure

**Return Value**

negative value reports error.

src/sja1000p.c

## sja1000p\_remote\_request

### Name

`sja1000p_remote_request` — configures message object and asks for RTR message

### Synopsis

```
int sja1000p_remote_request (struct chip_t * chip, struct msgobj_t * obj);
```

### Arguments

*chip*

pointer to chip state structure

*obj*

pointer to message object structure

### Return Value

negative value reports error.

**File**

src/sja1000p.c

## sja1000p\_standard\_mask

### Name

`sja1000p_standard_mask` — setup of mask for message filtering

### Synopsis

```
int sja1000p_standard_mask (struct chip_t * chip, unsigned short code,  
                           unsigned short mask);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*chip*

pointer to chip state structure

*code*

can message acceptance code

*mask*

can message acceptance mask

## Return Value

negative value reports error.

## File

src/sja1000p.c

# sja1000p\_clear\_objects

## Name

`sja1000p_clear_objects` — clears state of all message object residing in chip

## Synopsis

```
int sja1000p_clear_objects (struct chip_t * chip);
```

## Arguments

*chip*

pointer to chip state structure

## Return Value

negative value reports error.

## File

src/sja1000p.c

**Name**

`sja1000p_config_irqs` — tunes chip hardware interrupt delivery

**Synopsis**

```
int sja1000p_config_irqs (struct chip_t * chip, short irqs);
```

**Arguments**

*chip*

pointer to chip state structure

*irqs*

requested chip IRQ configuration

**Return Value**

negative value reports error.

**File**

`src/sja1000p.c`

**sja1000p\_irq\_write\_handler****Name**

`sja1000p_irq_write_handler` — part of ISR code responsible for transmit events

**Synopsis**

```
void sja1000p_irq_write_handler (struct chip_t * chip, struct msgobj_t * obj);
```

**Arguments**

*chip*

pointer to chip state structure

pointer to attached queue description

## Description

The main purpose of this function is to read message from attached queues and transfer message contents into CAN controller chip. This subroutine is called by `sja1000p_irq_write_handler` for transmit events.

## File

src/sja1000p.c

## **sja1000p\_irq\_handler**

### Name

`sja1000p_irq_handler` — interrupt service routine

### Synopsis

```
irqreturn_t sja1000p_irq_handler (int irq, void * dev_id, struct  
pt_regs * regs);
```

### Arguments

*irq*

interrupt vector number, this value is system specific

*dev\_id*

driver private pointer registered at time of `request_irq` call. The CAN driver uses this pointer to store relationship of interrupt to chip state structure - `struct chip_t`

*regs*

system dependent value pointing to registers stored in exception frame

## Description

Interrupt handler is activated when state of CAN controller chip changes, there is message to be read or there is more space for new messages or error occurs. The receive events results in reading of the message from CAN controller chip and distribution of message through attached message queues.

## File

src/sja1000p.c

**Name**

`sja1000p_wakeup_tx` — wakeups TX processing

**Synopsis**

```
int sja1000p_wakeup_tx (struct chip_t * chip, struct msgobj_t * obj);
```

**Arguments**

*chip*

pointer to chip state structure

*obj*

pointer to message object structure

**Return Value**

negative value reports error.

**File**

`src/sja1000p.c`

## CAN Queues Structures and Functions

**struct canque\_slot\_t****Name**

`struct canque_slot_t` — one CAN message slot in the CAN FIFO queue

**Synopsis**

```
struct canque_slot_t {
    struct canque_slot_t * next;
    unsigned long slot_flags;
    struct canmsg_t msg;
};
```

**Members**

*next*

pointer to the next/younger slot

slot\_flags

*Chapter 3. Driver Internal Documentation*

space for flags and optional command describing action associated with slot data

msg

space for one CAN message

## Description

This structure is used to store CAN messages in the CAN FIFO queue.

## **struct canque\_fifo\_t**

### Name

`struct canque_fifo_t` — CAN FIFO queue representation

### Synopsis

```
struct canque_fifo_t {
    unsigned long fifo_flags;
    unsigned long error_code;
    struct canque_slot_t * head;
    struct canque_slot_t ** tail;
    struct canque_slot_t * flist;
    struct canque_slot_t * entry;
    spinlock_t fifo_lock;
};
```

### Members

fifo\_flags

this field holds global flags describing state of the FIFO. `CAN_FIFOF_ERROR` is set when some error condition occurs. `CAN_FIFOF_ERR2BLOCK` defines, that error should lead to the FIFO block state. `CAN_FIFOF_BLOCK` state blocks insertion of the next messages. `CAN_FIFOF_OVERRUN` attempt to acquire new slot, when FIFO is full. `CAN_FIFOF_FULL` indicates FIFO full state. `CAN_FIFOF_EMPTY` indicates no allocated slot in the FIFO. `CAN_FIFOF_DEAD` condition indication. Used when FIFO is being destroyed.

error\_code

further description of error condition

head

pointer to the FIFO head, oldest slot

tail

pointer to the location, where pointer to newly inserted slot should be added

flist

pointer to list of the free slots associated with queue

entry

pointer to the memory allocated for the list slots.

the lock to ensure atomicity of slot manipulation operations.

## Description

This structure represents CAN FIFO queue. It is implemented as a single linked list of slots prepared for processing. The empty slots are stored in single linked list (*flist*).

## canque\_fifo\_get\_inslot

### Name

canque\_fifo\_get\_inslot — allocate slot for the input of one CAN message

### Synopsis

```
int canque_fifo_get_inslot (struct canque_fifo_t * fifo, struct  
canque_slot_t ** slotp, int cmd);
```

### Arguments

*fifo*

pointer to the FIFO structure

*slotp*

pointer to location to store pointer to the allocated slot.

*cmd*

optional command associated with allocated slot.

### Return Value

The function returns negative value if there is no free slot in the FIFO queue.

## canque\_fifo\_put\_inslot

### Name

canque\_fifo\_put\_inslot — releases slot to further processing

### Synopsis

```
int canque_fifo_put_inslot (struct canque_fifo_t * fifo, struct  
canque_slot_t * slot);
```

## Arguments

*fifo*

pointer to the FIFO structure

*slot*

pointer to the slot previously acquired by `canque_fifo_get_inslot`.

## Return Value

The nonzero return value indicates, that the queue was empty before call to the function. The caller should wake-up output side of the queue.

## **canque\_fifo\_abort\_inslot**

### Name

`canque_fifo_abort_inslot` — release and abort slot

### Synopsis

```
int canque_fifo_abort_inslot (struct canque_fifo_t * fifo, struct
canque_slot_t * slot);
```

## Arguments

*fifo*

pointer to the FIFO structure

*slot*

pointer to the slot previously acquired by `canque_fifo_get_inslot`.

## Return Value

The nonzero value indicates, that fifo was full

## **canque\_fifo\_test\_outslot**

### Name

`canque_fifo_test_outslot` — test and get ready slot from the FIFO

```
int canque_fifo_test_outslot (struct canque_fifo_t * fifo, struct  
canque_slot_t ** slotp);
```

**Arguments**

*fifo*

pointer to the FIFO structure

*slotp*

pointer to location to store pointer to the oldest slot from the FIFO.

**Return Value**

The negative value indicates, that queue is empty. The positive or zero value represents command stored into slot by the call to the function `canque_fifo_get_inslot`. The successfully acquired FIFO output slot has to be released by the call `canque_fifo_free_outslot` or `canque_fifo_again_outslot`.

**canque\_fifo\_free\_outslot****Name**

`canque_fifo_free_outslot` — free processed FIFO slot

**Synopsis**

```
int canque_fifo_free_outslot (struct canque_fifo_t * fifo, struct  
canque_slot_t * slot);
```

**Arguments**

*fifo*

pointer to the FIFO structure

*slot*

pointer to the slot previously acquired by `canque_fifo_test_outslot`.

**Return Value**

The returned value informs about FIFO state change. The mask `CAN_FIFOF_FULL` indicates, that the FIFO was full before the function call. The mask `CAN_FIFOF_EMPTY` informs, that last ready slot has been processed.

**Name**

canque\_fifo\_again\_outslot — interrupt and postpone processing of the slot

**Synopsis**

```
int canque_fifo_again_outslot (struct canque_fifo_t * fifo, struct  
canque_slot_t * slot);
```

**Arguments**

*fifo*

pointer to the FIFO structure

*slot*

pointer to the slot previously acquired by `canque_fifo_test_outslot`.

**Return Value**

The function cannot fail..

**struct canque\_edge\_t****Name**

`struct canque_edge_t` — CAN message delivery subsystem graph edge

**Synopsis**

```
struct canque_edge_t {  
    struct canque_fifo_t fifo;  
    unsigned long filtid;  
    unsigned long filtmask;  
    struct list_head inpeers;  
    struct list_head outpeers;  
    struct canque_ends_t * inends;  
    struct canque_ends_t * outends;  
    atomic_t edge_used;  
    int edge_prio;  
    int edge_num;  
};
```

**Members**

*fifo*

place where primitive `struct canque_fifo_t` FIFO is located.

the possible CAN message identifiers filter.

#### filtmask

the filter mask, the comparison considers only *filtid* bits corresponding to set bits in the *filtmask* field.

#### inpeers

the lists of all peer FIFOs connected by their input side (*inends*) to the same terminal (*struct canque\_ends\_t*).

#### outpeers

the lists of all peer FIFOs connected by their output side (*outends*) to the same terminal (*struct canque\_ends\_t*).

#### inends

the pointer to the FIFO input side terminal (*struct canque\_ends\_t*).

#### outends

the pointer to the FIFO output side terminal (*struct canque\_ends\_t*).

#### edge\_used

the atomic usage counter, mainly used for safe destruction of the edge.

#### edge\_prio

the assigned queue priority from the range 0 to CANQUEUE\_PRIO\_NR-1

#### edge\_num

edge sequential number intended for debugging purposes only

## Description

This structure represents one direction connection from messages source (*inends*) to message consumer (*outends*) fifo ends hub. The edge contains *&struct canque\_fifo\_t* for message fifo implementation.

## struct canque\_ends\_t

### Name

`struct canque_ends_t` — CAN message delivery subsystem graph vertex (FIFO ends)

### Synopsis

```
struct canque_ends_t {
    struct list_head * active;
    struct list_head idle;
    struct list_head inlist;
    spinlock_t ends_lock;
    void (* notify (struct canque_ends_t *qends, struct canque_edge_t *qedge, int what));
    void * context;
    union endinfo;
};
```

active

the array of the lists of active edges directed to the ends structure with ready messages. The array is indexed by the edges priorities.

idle

the list of the edges directed to the ends structure with empty FIFOs.

inlist

the list of outgoing edges input sides.

ends\_lock

the lock synchronizing operations between threads accessing same ends structure.

notify

pointer to notify procedure. The next state changes are notified. CANQUEUE\_NOTIFY\_EMPTY (out->in call) - all slots are processed by FIFO out side. CANQUEUE\_NOTIFY\_SPACE (out->in call) - full state negated => there is space for new message. CANQUEUE\_NOTIFY\_PROC (in->out call) - empty state negated => out side is requested to process slots. CANQUEUE\_NOTIFY\_NOUSR (both) - notify, that the last user has released the edge usage called with some lock to prevent edge disappear. CANQUEUE\_NOTIFY\_DEAD (both) - edge is in progress of deletion. CANQUEUE\_NOTIFY\_ATTACH (both) - new edge has been attached to end. CANQUEUE\_NOTIFY\_FILTER (out->in call) - edge filter rules changed CANQUEUE\_NOTIFY\_ERROR (out->in call) - error in messages processing.

context

space to store ends user specific information

endinfo

space to store some other ends usage specific informations mainly for waking-up by the notify calls.

## Description

Structure represents place to connect edges to for CAN communication entity. The zero, one or more incoming and outgoing edges can be connected to this structure.

## canque\_notify\_inends

### Name

canque\_notify\_inends — request to send notification to the input ends

### Synopsis

```
void canque_notify_inends (struct canque_edge_t * qedge, int what);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*qedge*

pointer to the edge structure

*what*

notification type

## canque\_notify\_outends

### Name

canque\_notify\_outends — request to send notification to the output ends

### Synopsis

```
void canque_notify_outends (struct canque_edge_t * qedge, int what);
```

## Arguments

*qedge*

pointer to the edge structure

*what*

notification type

## canque\_notify\_bothends

### Name

canque\_notify\_bothends — request to send notification to the both ends

### Synopsis

```
void canque_notify_bothends (struct canque_edge_t * qedge, int what);
```

## Arguments

*qedge*

pointer to the edge structure

notification type

## canque\_activate\_edge

### Name

canque\_activate\_edge — mark output end of the edge as active

### Synopsis

```
void canque_activate_edge (struct canque_ends_t * inends, struct
                           canque_edge_t * qedge);
```

### Arguments

*inends*

input side of the edge

*qedge*

pointer to the edge structure

### Description

Function call moves output side of the edge from idle onto active edges list.

## canque\_filtid2internal

### Name

canque\_filtid2internal — converts message ID and filter flags into internal format

### Synopsis

```
unsigned int canque_filtid2internal (unsigned long id, int filtflags);
```

### Arguments

*id*

CAN message 11 or 29 bit identifier

CAN message flags

## Description

This function maps message ID and `MSG_RTR`, `MSG_EXT` and `MSG_LOCAL` into one 32 bit number

# **canque\_fifo\_flush\_slots**

## Name

`canque_fifo_flush_slots` — free all ready slots from the FIFO

## Synopsis

```
int canque_fifo_flush_slots (struct canque_fifo_t * fifo);
```

## Arguments

*fifo*

pointer to the FIFO structure

## Description

The caller should be prepared to handle situations, when some slots are held by input or output side slots processing. These slots cannot be flushed or their processing interrupted.

## Return Value

The nonzero value indicates, that queue has not been empty before the function call.

# **canque\_fifo\_init\_slots**

## Name

`canque_fifo_init_slots` — initialize one CAN FIFO

## Synopsis

```
int canque_fifo_init_slots (struct canque_fifo_t * fifo, int slotsnr);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*fifo*

pointer to the FIFO structure

*slotsnr*

number of requested slots

## Return Value

The negative value indicates, that there is no memory to allocate space for the requested number of the slots.

## canque\_fifo\_done

### Name

canque\_fifo\_done — frees slots allocated for CAN FIFO

### Synopsis

```
int canque_fifo_done (struct canque_fifo_t * fifo);
```

## Arguments

*fifo*

pointer to the FIFO structure

## canque\_get\_inslot

### Name

canque\_get\_inslot — finds one outgoing edge and allocates slot from it

### Synopsis

```
int canque_get_inslot (struct canque_ends_t * qends, struct  
canque_edge_t ** qedgep, struct canque_slot_t ** slotp, int cmd);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*qends*

ends structure belonging to calling communication object

*qedgep*

place to store pointer to found edge

*slotp*

place to store pointer to allocated slot

*cmd*

command type for slot

## Description

Function looks for the first non-blocked outgoing edge in *qends* structure and tries to allocate slot from it.

## Return Value

If there is no usable edge or there is no free slot in edge negative value is returned.

## canque\_get\_inslot4id

### Name

canque\_get\_inslot4id — finds best outgoing edge and slot for given ID

### Synopsis

```
int canque_get_inslot4id (struct canque_ends_t * qends, struct  
canque_edge_t ** qedgep, struct canque_slot_t ** slotp, int cmd,  
unsigned long id, int prio);
```

## Arguments

*qends*

ends structure belonging to calling communication object

*qedgep*

place to store pointer to found edge

*slotp*

place to store pointer to allocated slot

*cmd*

command type for slot

communication ID of message to send into edge

optional priority of message

## Description

Function looks for the non-blocked outgoing edge accepting messages with given ID. If edge is found, slot is allocated from that edge. The edges with non-zero mask are preferred over edges open to all messages. If more edges with mask accepts given message ID, the edge with highest priority below or equal to required priority is selected.

## Return Value

If there is no usable edge or there is no free slot in edge negative value is returned.

# canque\_put\_inslot

## Name

canque\_put\_inslot — schedules filled slot for processing

## Synopsis

```
int canque_put_inslot (struct canque_ends_t * qends, struct  
canque_edge_t * qedge, struct canque_slot_t * slot);
```

## Arguments

ends structure belonging to calling communication object

edge slot belong to

pointer to the prepared slot

## Description

Puts slot previously acquired by `canque_get_inslot` or `canque_get_inslot4id` function call into FIFO queue and activates edge processing if needed.

## Return Value

## Chapter 3. Driver Internal Documentation

Positive value informs, that activation of output end has been necessary

## canque\_abort\_inslot

### Name

canque\_abort\_inslot — aborts preparation of the message in the slot

### Synopsis

```
int canque_abort_inslot (struct canque_ends_t * qends, struct  
canque_edge_t * qedge, struct canque_slot_t * slot);
```

### Arguments

*qends*

ends structure belonging to calling communication object

*qedge*

edge slot belong to

*slot*

pointer to the previously allocated slot

### Description

Frees slot previously acquired by `canque_get_inslot` or `canque_get_inslot4id` function call. Used when message copying into slot fails.

## Return Value

Positive value informs, that queue full state has been negated.

## canque\_filter\_msg2edges

### Name

canque\_filter\_msg2edges — sends message into all edges which accept its ID

### Synopsis

```
int canque_filter_msg2edges (struct canque_ends_t * qends, struct  
canmsg_t * msg);
```

## Arguments

*qends*

ends structure belonging to calling communication object

*msg*

pointer to CAN message

## Description

Sends message to all outgoing edges connected to the given ends, which accepts message communication ID.

## Return Value

Returns number of edges message has been send to

## canque\_test\_outslot

### Name

canque\_test\_outslot — test and retrieve ready slot for given ends

### Synopsis

```
int canque_test_outslot (struct canque_ends_t * qends, struct
canque_edge_t ** qedgep, struct canque_slot_t ** slotp);
```

## Arguments

*qends*

ends structure belonging to calling communication object

*qedgep*

place to store pointer to found edge

*slotp*

place to store pointer to received slot

## Description

Function takes highest priority active incoming edge and retrieves oldest ready slot from it.

## Return Value

## Chapter 3. Driver Internal Documentation

Negative value informs, that there is no ready output slot for given ends. Positive value is equal to the command slot has been allocated by the input side.

## canque\_free\_outslot

### Name

canque\_free\_outslot — frees processed output slot

### Synopsis

```
int canque_free_outslot (struct canque_ends_t * qends, struct  
canque_edge_t * qedge, struct canque_slot_t * slot);
```

### Arguments

*qends*

ends structure belonging to calling communication object

*qedge*

edge slot belong to

*slot*

pointer to the processed slot

### Description

Function releases processed slot previously acquired by `canque_test_outslot` function call.

## Return Value

Return value informs if input side has been notified to know about change of edge state

## canque\_again\_outslot

### Name

canque\_again\_outslot — reschedule output slot to process it again later

## Synopsis

## Chapter 3. Driver Internal Documentation

```
int canque_again_outslot (struct canque_ends_t * qends, struct  
canque_edge_t * qedge, struct canque_slot_t * slot);
```

## Arguments

*qends*

ends structure belonging to calling communication object

*qedge*

edge slot belong to

*slot*

pointer to the slot for re-processing

## Description

Function reschedules slot previously acquired by `canque_test_outslot` function call for second time processing.

## Return Value

Function cannot fail.

## canque\_set\_filt

### Name

`canque_set_filt` — sets filter for specified edge

## Synopsis

```
int canque_set_filt (struct canque_edge_t * qedge, unsigned long  
filtid, unsigned long filtmask, int filtflags);
```

## Arguments

*qedge*

pointer to the edge

*filtid*

ID to set for the edge

*filtmask*

mask used for ID match check

required filer flags

## Return Value

Negative value is returned if edge is in the process of delete.

# **canque\_flush**

## Name

canque\_flush — flushes all ready slots in the edge

## Synopsis

```
int canque_flush (struct canque_edge_t * qedge);
```

## Arguments

*qedge*

pointer to the edge

## Description

Tries to flush all allocated slots from the edge, but there could exist some slots associated to edge which are processed by input or output side and cannot be flushed at this moment.

## Return Value

The nonzero value indicates, that queue has not been empty before the function call.

# **canqueue\_ends\_init\_gen**

## Name

canqueue\_ends\_init\_gen — subsystem independent routine to initialize ends state

## Synopsis

```
int canqueue_ends_init_gen (struct canque_ends_t * qends);
```

*qends*

pointer to the ends structure

**Return Value**

Cannot fail.

**canqueue\_notify\_kern****Name**

canqueue\_notify\_kern — notification callback handler for Linux userspace clients

**Synopsis**

```
void canqueue_notify_kern (struct canque_ends_t * qends, struct
canque_edge_t * qedge, int what);
```

**Arguments***qends*

pointer to the callback side ends structure

*qedge*

edge which invoked notification

*what*

notification type

**canqueue\_ends\_init\_kern****Name**

canqueue\_ends\_init\_kern — Linux userspace clients specific ends initialization

**Synopsis**

```
int canqueue_ends_init_kern (struct canque_ends_t * qends);
```

*qends*

pointer to the callback side ends structure

## canque\_get\_inslot4id\_wait\_kern

### Name

`canque_get_inslot4id_wait_kern` — find or wait for best outgoing edge and slot for given ID

### Synopsis

```
int canque_get_inslot4id_wait_kern (struct canque_ends_t * qends,  
 struct canque_edge_t ** qedgep, struct canque_slot_t ** slotp, int  
 cmd, unsigned long id, int prio);
```

### Arguments

*qends*

ends structure belonging to calling communication object

*qedgep*

place to store pointer to found edge

*slotp*

place to store pointer to allocated slot

*cmd*

command type for slot

*id*

communication ID of message to send into edge

*prio*

optional priority of message

### Description

Same as `canque_get_inslot4id`, except, that it waits for free slot in case, that queue is full. Function is specific for Linux userspace clients.

### Return Value

If there is no usable edge negative value is returned.

### Name

canque\_get\_outslot\_wait\_kern — receive or wait for ready slot for given ends

### Synopsis

```
int canque_get_outslot_wait_kern (struct canque_ends_t * qends, struct  
canque_edge_t ** qedgep, struct canque_slot_t ** slotp);
```

### Arguments

*qends*

ends structure belonging to calling communication object

*qedgep*

place to store pointer to found edge

*slotp*

place to store pointer to received slot

### Description

The same as `canque_test_outslot`, except it waits in the case, that there is no ready slot for given ends. Function is specific for Linux userspace clients.

### Return Value

Negative value informs, that there is no ready output slot for given ends. Positive value is equal to the command slot has been allocated by the input side.

## **canque\_sync\_wait\_kern**

### Name

canque\_sync\_wait\_kern — wait for all slots processing

### Synopsis

```
int canque_sync_wait_kern (struct canque_ends_t * qends, struct  
canque_edge_t * qedge);
```

## Arguments

## Chapter 3. Driver Internal Documentation

*qends*

ends structure belonging to calling communication object

*qedge*

pointer to edge

## Description

Functions waits for ends transition into empty state.

## Return Value

Positive value indicates, that edge empty state has been reached. Negative or zero value informs about interrupted wait or other problem.

## **canque\_new\_edge\_kern**

### Name

`canque_new_edge_kern` — allocate new edge structure in the Linux kernel context

### Synopsis

```
struct canque_edge_t * canque_new_edge_kern (int slotsnr);
```

## Arguments

*slotsnr*

required number of slots in the newly allocated edge structure

## Return Value

Returns pointer to allocated slot structure or `NULL` if there is not enough memory to process operation.

## **canqueue\_connect\_edge**

### Name

`canqueue_connect_edge` — connect edge between two communication entities

## Synopsis

## Chapter 3. Driver Internal Documentation

```
int canqueue_connect_edge (struct canque_edge_t * qedge, struct  
canque_ends_t * inends, struct canque_ends_t * outends);
```

## Arguments

*qedge*

pointer to edge

*inends*

pointer to ends the input of the edge should be connected to

*outends*

pointer to ends the output of the edge should be connected to

## Return Value

Negative value informs about failed operation.

## canqueue\_disconnect\_edge

### Name

canqueue\_disconnect\_edge — disconnect edge from communicating entities

## Synopsis

```
int canqueue_disconnect_edge (struct canque_edge_t * qedge);
```

## Arguments

*qedge*

pointer to edge

## Return Value

Negative value means, that edge is used and cannot be disconnected. Operation has to be delayed.

### Name

canqueue\_disconnect\_edge\_kern — disconnect edge from communicating entities with wait

### Synopsis

```
int canqueue_disconnect_edge_kern (struct canque_ends_t * qends, struct  
canque_edge_t * qedge);
```

### Arguments

*qends*

ends structure belonging to calling communication object

*qedge*

pointer to edge

### Description

Same as `canqueue_disconnect_edge`, but tries to wait for state with zero use counter.

### Return Value

Negative value means, that edge is used and cannot be disconnected yet. Operation has to be delayed.

## **canqueue\_ends\_done\_kern**

### Name

canqueue\_ends\_done\_kern — finalizing of the ends structure for Linux kernel clients

### Synopsis

```
int canqueue_ends_done_kern (struct canque_ends_t * qends, int sync);
```

### Arguments

*qends*

pointer to ends structure

flag indicating, that user wants to wait for processing of all remaining messages

## **Return Value**

Function should be designed such way to not fail.